

1.

- (a) Let \mathcal{S} be the statement

For all integers n , if n is even then $3n - 11$ is odd.

Is \mathcal{S} true? Give a proof or counterexample.

- (b) Write out the *contrapositive* of statement \mathcal{S} , and give a proof or disproof.
 (c) Write out the *converse* of statement \mathcal{S} , and give a proof or disproof.
 (d) Prove or disprove the statement

For all integers n , if n is odd then $2n - 11$ is even.

Then write out the converse of this statement and prove or disprove it.

- (a) \mathcal{S} is **true**. Here is a proof.

Let n be an arbitrary even integer. This means that $n = 2k$ for some integer k . Then

$$3n - 11 = 3(2k) - 11 = 6k - 11 = 2(3k - 6) + 1$$

where $3k - 6$ is an integer. Therefore $3n - 11$ is odd by the definition of odd.

- (b) The contrapositive of \mathcal{S} is:

For all integers n , if $3n - 11$ is not odd then n is not even,

which could also be written (using a result on page 159 of the text)

For all integers n , if $3n - 11$ is even then n is odd.

It is **true**, because it is equivalent to the original statement \mathcal{S} which is true.

- (c) The converse of \mathcal{S} is

For all integers n , if $3n - 11$ is odd then n is even.

This statement is **true**. Here is a proof.

Assume that $3n - 11$ is odd, where n is an integer. This means that $3n - 11 = 2k + 1$ for some integer k . We can rewrite this equation as $n = 2k + 12 - 2n = 2(k + 6 - n)$, where $k + 6 - n$ is an integer since k and n are integers. Therefore n equals 2 times an integer, so n is even.

Note. The converse could also be proven by writing its contrapositive

For all integers n , if n is not even then $3n - 11$ is not odd

in the form

For all integers n , if n is odd then $3n - 11$ is even

and proving this.

- (d) This statement is **false**. A counterexample is $n = 1$. For then n is odd, but $2n - 11 = 2 - 11 = -9$ is not even.

The converse of this statement is

For all integers n , if $2n - 11$ is even then n is odd.

This statement is **true** vacuously. For every integer n , $2n - 11 = 2(n - 6) + 1$ where $n - 6$ is an integer, thus $2n - 11$ is odd and so cannot be even. Since the “if” part of the conditional never holds, the statement is true vacuously.

2. Prove or disprove the following statements:

- (a) There exists a prime number a such that $a + 271$ is prime.
(b) There exists a prime number a such that $a + 271$ is composite.
(c) There exists a composite number a such that $a + 271$ is prime.
(d) There exists a composite number a such that $a + 271$ is composite.
(e) Choose one of statements (a) to (d) (your choice), replace 271 with your U of C ID number, and prove or disprove the resulting statement.

- (a) This statement is **false**. Here is a proof.

Assume a is a prime number. We have two cases.

Case (i): $a = 2$. Then $a + 271 = 2 + 271 = 273 = 3 \cdot 91$, so $a + 271$ is not prime.

Case (ii): $a > 2$. Then a must be odd, so $a = 2k + 1$ for some integer k . Then $a + 271 = 2k + 1 + 271 = 2k + 272 = 2(k + 136)$, where $k + 136$ is an integer. Therefore $a + 271$ is not prime.

In neither case can we get that $a + 271$ is prime, so the statement is false.

- (b) This statement is **true**. An example is $a = 3$. Then a is prime and $a + 271 = 274 = 2 \cdot 137$ is composite.
(c) This statement is **true**. An example is $a = 6$. Then a is composite and $a + 271 = 277$ is prime (it turns out).

Note. An alternate proof would go like this: since there are infinitely many primes (Theorem 3.7.4 of the text), there must be a prime $p \geq 275$. Then p is odd, so $p - 271$ must be even (prove it), and $p - 271 \geq 4$, so $p - 271$ is composite. Put $a = p - 271$; then a is composite and $a + 271 = p$ is prime.

- (d) This statement is **true**. An example is $a = 9$. Then a is composite and $a + 271 = 280$ is composite too.
(e) Regardless of what your ID number is, probably (d) is the easiest statement to prove. Let's do it for the hypothetical ID number 123456. Choosing $a = 4$, we get that a is composite and that $a + 123456 = 123460$ is also composite.

3. *Note:* \mathbf{Z} denotes the set of all integers, and \mathbf{Z}^+ denotes the set of all positive integers.

(a) Prove the following statements:

- (i) $\exists a \in \mathbf{Z}$ so that $\forall b \in \mathbf{Z}, (a - b)|(a + b)$.
- (ii) $\forall a \in \mathbf{Z}^+ \exists b \in \mathbf{Z}^+$ so that $(a - b)|(a + b)$.
- (iii) $\forall a \in \mathbf{Z}^+ \exists b \in \mathbf{Z}^+$ so that $(a + b)|(a - b)$.

(b) Write out the *negation* of the following statement:

$$\forall a, b \in \mathbf{Z}^+, \text{ if } a|2 \text{ and } b|3 \text{ then } (a + b)|5.$$

Then show that the negation is true, so that the original statement is false.

(c) Prove the following statement:

$$\exists N \in \mathbf{Z}^+ \text{ so that } \forall a, b \in \mathbf{Z}^+, \text{ if } a|2 \text{ and } b|3 \text{ then } (a + b)|N.$$

(a) (i) Choose $a = 0$. Then the statement to be proved is: $\forall b \in \mathbf{Z}, (-b)|b$. To prove this, let b be an arbitrary integer. Then $b = (-b)(-1)$ where -1 is an integer, so $(-b)|b$.

(ii) Let a be an arbitrary positive integer. We need to find a positive integer b (maybe depending on a) so that $(a - b)|(a + b)$. Choose $b = a + 1$, which is a positive integer. Then $a - b = -1$ and $a + b = 2a + 1$, so we need to show that $(-1)|(2a + 1)$. But this is clear, since $2a + 1 = (-1)(-2a - 1)$ where $-2a - 1$ is an integer.

(iii) Let a be an arbitrary positive integer. We need to find a positive integer b (maybe depending on a) so that $(a + b)|(a - b)$. Choose $b = a$, which is a positive integer. Then $a + b = 2a$ and $a - b = 0$, so we need to show that $(2a)|0$. But this is clear, since $0 = 0 \cdot 2a$.

(b) The negation is:

$$\exists a, b \in \mathbf{Z}^+ \text{ so that } a|2 \text{ and } b|3 \text{ but } (a + b) \nmid 5.$$

This statement is true. For example we can choose $a = 1$ and $b = 1$; then $a|2$ and $b|3$ are both true, but $a + b = 2$, and $2 \nmid 5$.

(c) For $a|2$ we need either $a = 1$ or $a = 2$, and for $b|3$ we need either $b = 1$ or $b = 3$. Thus we will need N to satisfy all of the following:

- $(1 + 1)|N$, which says $2|N$;
- $(2 + 1)|N$, which says $3|N$;
- $(1 + 3)|N$, which says $4|N$;
- $(2 + 3)|N$, which says $5|N$.

So for example, $N = 2 \cdot 3 \cdot 4 \cdot 5 = \mathbf{120}$ will work. Actually $N = 3 \cdot 4 \cdot 5 = \mathbf{60}$ will work too, and this is the smallest value of N which will work.

MATH 271 ASSIGNMENT 2 SOLUTIONS

1. (a) Find all positive integers a so that $\lfloor a/271 \rfloor = 10$. How many such integers are there?
 - (b) Find all positive integers a so that $\lfloor 271/a \rfloor = 10$.
 - (c) Find all positive integers a so that $\lceil 271/a \rceil = 10$.
 - (d) Prove or disprove: $\forall n \in \mathbf{Z}$, the equations $\lfloor 271/x \rfloor = n$ and $\lceil 271/x \rceil = n$ have the *same number* of integer solutions x .
 - (e) Prove or disprove: $\exists n \in \mathbf{Z}$ so that $\forall a \in \mathbf{Z}$, $\lfloor 271/a \rfloor \neq n$.
- (a) For $\lfloor a/271 \rfloor = 10$ to be true we would need $10 \leq a/271 < 11$, or $2710 \leq a < 271 \cdot 11 = 2981$. So the values of a are 2710, 2711, 2712, \dots , 2980, a total of 271 integers.
- (b) Now we will need $10 \leq 271/a < 11$, or $10a \leq 271 < 11a$. This means $a \leq 271/10$ and $a > 271/11$, in other words $24.6 < a \leq 27.1$. So the allowed values of a are 25, 26, 27.
- (c) Similarly, this time we will need $9 < 271/a \leq 10$, or $9a < 271 \leq 10a$. This means $a < 271/9$ and $a \geq 271/10$, in other words $27.1 \leq a < 30.1$. So the allowed values of a are 28, 29, 30.
- (d) Despite the “evidence” from parts (b) and (c) (where there were 3 solutions each time), this statement is *false*. One counterexample is $n = 11$, as the only solutions for $\lfloor 271/x \rfloor = 11$ are $x = 23$ and 24, while $\lceil 271/x \rceil = 11$ has the three solutions $x = 25, 26$ and 27. Another counterexample is $n = 8$, since $\lfloor 271/x \rfloor = 8$ has three solutions $x = 31, 32$ and 33 while $\lceil 271/x \rceil = 8$ has the five solutions $x = 34$ to 38.
- An interesting counterexample is $n = 1$. Notice that the equation $\lceil 271/x \rceil = 1$ means $0 < 271/x \leq 1$, which is satisfied for *every* integer x greater than or equal to 271. So there are infinitely many solutions. But the equation $\lfloor 271/x \rfloor = 1$ means $1 \leq 271/x < 2$, and this inequality is satisfied only for the integers $x = 136, 137, \dots, 271$.
- (e) This statement is *true*, and there are lots of integers n satisfying the condition. For example, any $n > 271$ will work, because $\lfloor 271/a \rfloor > 271$ is impossible for an integer a . *Note.* Can you find the *smallest* positive integer n for which this statement is true? If you think you have an answer to this question, talk to your professor or TA.

2. Let

$$S_n = \frac{3}{1 \cdot 2} - \frac{5}{2 \cdot 3} + \frac{7}{3 \cdot 4} - \frac{9}{4 \cdot 5} + \cdots - \frac{4n-3}{(2n-2)(2n-1)} + \frac{4n-1}{(2n-1)2n},$$

where the signs alternate.

- (a) Calculate and simplify S_1 , S_2 and S_3 .
- (b) Use part (a) (and more calculations if you need them) to guess a simple formula for S_n .
- (c) Prove your formula for all positive integers n *using mathematical induction*.
- (d) Give another proof of your formula for all positive integers n *using telescoping*. (See example 4.1.10 on page 205 of the text.)

(a) We get

$$S_1 = \frac{3}{1 \cdot 2} = \frac{3}{2}, \quad S_2 = \frac{3}{1 \cdot 2} - \frac{5}{2 \cdot 3} + \frac{7}{3 \cdot 4} = \frac{3}{2} - \frac{5}{6} + \frac{7}{12} = \frac{18 - 10 + 7}{12} = \frac{15}{12} = \frac{5}{4},$$

and (using our calculation for S_2)

$$S_3 = \frac{3}{1 \cdot 2} - \frac{5}{2 \cdot 3} + \frac{7}{3 \cdot 4} - \frac{9}{4 \cdot 5} + \frac{11}{5 \cdot 6} = \frac{5}{4} - \frac{9}{20} + \frac{11}{30} = \frac{75 - 27 + 22}{60} = \frac{70}{60} = \frac{7}{6}.$$

(b) From the values in (a) we guess that $S_n = \frac{2n+1}{2n}$.

(c) *Basis step.* We need to prove that $S_1 = \frac{2 \cdot 1 + 1}{2 \cdot 1}$, which is true since both are $3/2$.

Induction step. Assume that $S_k = \frac{2k+1}{2k}$ for some integer $k \geq 1$. We want to prove that $S_{k+1} = \frac{2(k+1)+1}{2(k+1)}$, which is the same as $S_{k+1} = \frac{2k+3}{2(k+1)}$. Well,

$$\begin{aligned} S_{k+1} &= \frac{3}{1 \cdot 2} - \frac{5}{2 \cdot 3} + \cdots + \frac{4k-1}{(2k-1)2k} - \frac{4k+1}{2k(2k+1)} + \frac{4k+3}{(2k+1)(2k+2)} \\ &= S_k - \frac{4k+1}{2k(2k+1)} + \frac{4k+3}{(2k+1)(2k+2)} \\ &= \frac{2k+1}{2k} - \frac{4k+1}{2k(2k+1)} + \frac{4k+3}{(2k+1)(2k+2)} \quad \text{by our assumption} \\ &= \frac{(2k+1)^2(2k+2) - (4k+1)(2k+2) + (4k+3)2k}{2k(2k+1)(2k+2)} \\ &= \frac{[4k^2 + 4k + 1 - (4k+1)](2k+2) + (8k^2 + 6k)}{2k(2k+1)(2k+2)} \\ &= \frac{4k^2(2k+2) + (8k^2 + 6k)}{2k(2k+1)(2k+2)} = \frac{8k^3 + 16k^2 + 6k}{2k(2k+1)(2k+2)} \\ &= \frac{2k(2k+1)(2k+3)}{2k(2k+1)(2k+2)} = \frac{2k+3}{2k+2}, \end{aligned}$$

which proves the induction step.

Therefore the statement is true for all integers $n \geq 1$.

(d) Notice that

$$\frac{3}{1 \cdot 2} = \frac{1}{1} + \frac{1}{2}, \quad \frac{5}{2 \cdot 3} = \frac{1}{2} + \frac{1}{3}, \quad \frac{7}{3 \cdot 4} = \frac{1}{3} + \frac{1}{4},$$

and in general

$$\frac{2k+1}{k(k+1)} = \frac{1}{k} + \frac{1}{k+1}$$

for any positive integer k . Thus

$$\begin{aligned}
S_n &= \frac{3}{1 \cdot 2} - \frac{5}{2 \cdot 3} + \frac{7}{3 \cdot 4} - \frac{9}{4 \cdot 5} + \cdots - \frac{4n-3}{(2n-2)(2n-1)} + \frac{4n-1}{(2n-1)2n} \\
&= \left(\frac{1}{1} + \frac{1}{2}\right) - \left(\frac{1}{2} + \frac{1}{3}\right) + \left(\frac{1}{3} + \frac{1}{4}\right) - \cdots - \left(\frac{1}{2n-2} + \frac{1}{2n-1}\right) + \left(\frac{1}{2n-1} + \frac{1}{2n}\right) \\
&= \frac{1}{1} + \frac{1}{2n} = \frac{2n+1}{2n},
\end{aligned}$$

so our guess is proved.

3. (a) Prove the following statement by contradiction: for all integers n , if $3|n$ then $3 \nmid (n+271)$.
- (b) Prove or disprove: for all integers n , if $3|n$ then $5 \nmid n$.
- (c) Prove by mathematical induction that $3|(2^n - (-1)^n)$ for all integers $n \geq 1$.

- (a) Assume that $3|n$ for some integer n . This means that $n = 3k$ for some integer k . We want to prove that $3 \nmid (n+271)$. To get a proof by contradiction, we assume that what we want to prove is false: namely we will assume that $3|(n+271)$. This means we are also assuming that $n+271 = 3\ell$ for some integer ℓ . Now our two assumptions tell us that

$$271 = (n+271) - n = 3\ell - 3k = 3(\ell - k),$$

where $\ell - k$ is an integer. Thus $3|271$, which however is false. Thus our assumption that $3|(n+271)$ must be false, so $3 \nmid (n+271)$.

- (b) This is *false*. A counterexample is $n = 15$, since $3|15$ but also $5|15$. Another counterexample is $n = 0$.
- (c) *Basis step.* We need to prove that $3|(2^1 - (-1)^1)$, which says $3|(2+1)$ or $3|3$. This is true.

Induction step. Assume that $3|(2^k - (-1)^k)$ for some integer $k \geq 1$. This means that $2^k - (-1)^k = 3\ell$ for some integer ℓ . We want to prove that $3|(2^{k+1} - (-1)^{k+1})$. Well,

$$\begin{aligned}
2^{k+1} - (-1)^{k+1} &= 2 \cdot 2^k - (-1) \cdot (-1)^k \\
&= 2(2^k - (-1)^k) + 2(-1)^k + (-1)^k \\
&= 2(3\ell) + 3(-1)^k \quad \text{by our assumption} \\
&= 3(2\ell + (-1)^k),
\end{aligned}$$

where $2\ell + (-1)^k$ is an integer. Thus $3|(2^{k+1} - (-1)^{k+1})$.

Therefore, by induction, $3|(2^n - (-1)^n)$ for all integers $n \geq 1$.

MATH 271 ASSIGNMENT 3 SOLUTIONS

1. (a) Prove **by induction** that, for all integers $n \geq 2$,

$$\frac{1^2}{2!} + \frac{2^2}{3!} + \frac{3^2}{4!} + \cdots + \frac{n^2}{(n+1)!} \leq 2 - \frac{2n}{(n+1)!}. \quad (1)$$

- (b) Prove that in fact inequality (1) holds for all integers $n \geq 1$.

- (c) Find the smallest real number A so that, for all integers $n \geq 1$,

$$\frac{1^2}{2!} + \frac{2^2}{3!} + \frac{3^2}{4!} + \cdots + \frac{n^2}{(n+1)!} \leq A - \frac{2n}{(n+1)!}.$$

- (a) *Basis step.* When $n = 2$ inequality (1) is

$$\frac{1^2}{2!} + \frac{2^2}{3!} \leq 2 - \frac{4}{3!}$$

which is

$$\frac{1}{2} + \frac{4}{6} \leq 2 - \frac{4}{6}, \quad \text{that is} \quad \frac{7}{6} \leq \frac{8}{6},$$

which is true.

Inductive step. Assume that inequality (1) holds for some integer $n = k$, where $k \geq 2$. We want to prove that inequality (1) holds for $n = k + 1$. So we are assuming that

$$\frac{1^2}{2!} + \frac{2^2}{3!} + \frac{3^2}{4!} + \cdots + \frac{k^2}{(k+1)!} \leq 2 - \frac{2k}{(k+1)!},$$

and we want to prove that

$$\frac{1^2}{2!} + \frac{2^2}{3!} + \frac{3^2}{4!} + \cdots + \frac{(k+1)^2}{(k+2)!} \leq 2 - \frac{2(k+1)}{(k+2)!}. \quad (2)$$

Well,

$$\begin{aligned} \frac{1^2}{2!} + \frac{2^2}{3!} + \cdots + \frac{(k+1)^2}{(k+2)!} &= \frac{1^2}{2!} + \frac{2^2}{3!} + \cdots + \frac{k^2}{(k+1)!} + \frac{(k+1)^2}{(k+2)!} \\ &\leq 2 - \frac{2k}{(k+1)!} + \frac{(k+1)^2}{(k+2)!} \quad \text{by our assumption} \\ &= 2 - \frac{2k(k+2) - (k+1)^2}{(k+2)!} \\ &= 2 - \frac{2k^2 + 4k - k^2 - 2k - 1}{(k+2)!} \\ &= 2 - \frac{k^2 + 2k - 1}{(k+2)!}. \end{aligned}$$

So in order to prove (2), we would like to prove that

$$2 - \frac{k^2 + 2k - 1}{(k + 2)!} \leq 2 - \frac{2(k + 1)}{(k + 2)!}.$$

This is equivalent successively to

$$-\frac{k^2 + 2k - 1}{(k + 2)!} \leq -\frac{2(k + 1)}{(k + 2)!},$$

$$\frac{k^2 + 2k - 1}{(k + 2)!} \geq \frac{2(k + 1)}{(k + 2)!},$$

and thus to

$$k^2 + 2k - 1 \geq 2k + 2, \quad \text{that is, } k^2 \geq 3,$$

which is true since $k \geq 2$. This finishes the proof of the inductive step. Thus inequality (1) holds for all integers $n \geq 2$.

(b) When $n = 1$, inequality (1) says

$$\frac{1^2}{2!} \leq 2 - \frac{2}{2!}$$

which is $1/2 \leq 1$, which is true. Since in part (a) we proved that inequality (1) holds for all integers $n \geq 2$, we now know it holds for all integers $n \geq 1$. Notice that, since the inductive step needed that $k \geq 2$, to prove inequality (1) for all $n \geq 1$ we need both cases $n = 1$ and $n = 2$ in the basis step.

(c) The inductive step in the proof in part (a) works just the same if the 2 right after the inequality sign is replaced with any number A . So the inequality in part (c) will hold for all integers $n \geq 1$ provided that it holds for $n = 1$ and $n = 2$, which is the basis step. When $n = 1$ the inequality in (c) says

$$\frac{1^2}{2!} \leq A - \frac{2}{2!}$$

which simplifies to $A \geq 3/2$. When $n = 2$ the inequality in (c) says

$$\frac{1^2}{2!} + \frac{2^2}{3!} \leq A - \frac{4}{3!}$$

which simplifies to $A \geq 1/2 + 4/6 + 4/6 = 11/6$. We need both of these to hold, so the smallest A that will work is $A = \mathbf{11/6}$.

2. You are given the following “while” loop:

[*Pre-condition*: m is a nonnegative integer, $a = 0$, $b = 1$, $c = 2$, $i = 0$.]

while ($i \neq m$)

1. $a := b$
2. $b := c$
3. $c := 2b - a$
4. $i := i + 1$

end while

[*Post-condition*: $c = m + 2$.]

Loop invariant: $I(n)$ is “ $a = n$, $b = n + 1$, $c = n + 2$, $i = n$ ”.

- (a) Prove the correctness of this loop with respect to the pre- and post-conditions.
 - (b) Suppose the “while” loop is as above, except that the pre-condition is replaced by: m is a nonnegative integer, $a = 1$, $b = 3$, $c = 5$, $i = 0$. Find a post-condition that gives the final value of c , and an appropriate loop invariant, and prove the correctness of this loop.
- (a) We first need to check that the loop invariant holds when $n = 0$. $I(0)$ says $a = 0$, $b = 1$, $c = 2$ and $i = 0$, and these are all true by the pre-conditions.

So now assume that the loop invariant $I(k)$ holds for some integer $k \geq 0$, $k < m$. We want to prove that $I(k + 1)$ holds, that is, that the loop invariant will still hold after one more pass through the loop. So we are assuming that $a = k$, $b = k + 1$, $c = k + 2$ and $i = k$, and we now go through the loop. Step 1 sets a equal to $b = k + 1$, then step 2 sets b equal to $c = k + 2$, then step 3 sets c equal to $2b - a = 2(k + 2) - (k + 1) = k + 3$, then step 4 sets i equal to $k + 1$. This means that $I(k + 1)$ is true, as required.

Finally the loop stops when $i = m$, and we need to check that at that point the post-condition is satisfied. When $i = m$ it means that the loop invariant $I(m)$ must hold, so from $I(m)$ we know that $c = m + 2$ as required.

- (b) If we set the variables to their pre-condition values of $a = 1$, $b = 3$, $c = 5$ and $i = 0$, and run through the loop, the new values we get are $a = 3$, $b = 5$, $c = 2(5) - 3 = 7$, and $i = 1$. From this (or by running through the loop once or twice more to collect more evidence) we can guess that the loop invariant we want will be

$$I(n) : a = 2n + 1, b = 2n + 3, c = 2n + 5, i = n,$$

and the post-condition value of c ought to be $c = 2m + 5$. This choice of $I(n)$ becomes $a = 1$, $b = 3$, $c = 5$ and $i = 0$ when $n = 0$, so the pre-condition is satisfied.

So now we assume that the new loop invariant $I(k)$ holds for some integer $k \geq 0$, $k < m$, and we want to prove that $I(k + 1)$ holds. So we are assuming that $a = 2k + 1$, $b = 2k + 3$, $c = 2k + 5$ and $i = k$, and we now go through the loop. Step 1 sets a equal to $b = 2k + 3 = 2(k + 1) + 1$, then step 2 sets b equal to $c = 2k + 5 = 2(k + 1) + 3$, then step 3 sets c equal to $2b - a = 2(2k + 5) - (2k + 3) = 2k + 7 = 2(k + 1) + 5$, then step 4 sets i equal to $k + 1$. This means that $I(k + 1)$ is true, as required.

Finally the loop stops when $i = m$, and we need to check that at that point the post-condition is satisfied. When $i = m$ it means that the loop invariant $I(m)$ must hold, so from $I(m)$ we know that $c = 2m + 5$ as required.

3. Prove or disprove each of the following six statements. Proofs should use the “element” methods given in Section 5.2. [Note: $\mathcal{P}(X)$ denotes the power set of the set X .]

- (a) For all sets A, B, C , $(A - B) \times C \subseteq (A \times C) - (B \times C)$.
- (b) For all sets A, B, C , $(A \times C) - (B \times C) \subseteq (A - B) \times C$.
- (c) For all sets A, B, C , $(A - B) \times C = (A \times C) - (B \times C)$.
- (d) For all sets A and B , $\mathcal{P}(A - B) \subseteq \mathcal{P}(A) - \mathcal{P}(B)$.
- (e) For all sets A and B , $\mathcal{P}(A) - \mathcal{P}(B) \subseteq \mathcal{P}(A - B)$.
- (f) For all sets A and B , $\mathcal{P}(A - B) = \mathcal{P}(A) - \mathcal{P}(B)$.

(a) This inequality is **true**. Here is a proof.

Let A, B, C be arbitrary sets. Note that the left side of this inequality is a Cartesian product, which means that its elements will be ordered pairs. So let (a, c) be an arbitrary element of $(A - B) \times C$. This means that $a \in A - B$ and $c \in C$. Since $a \in A - B$, this means that $a \in A$ and $a \notin B$. Since $a \in A$ and $c \in C$, we get that $(a, c) \in A \times C$. But since $a \notin B$, we know that (a, c) cannot be an element of $B \times C$. Since $(a, c) \in A \times C$ but $(a, c) \notin B \times C$, we know $(a, c) \in (A \times C) - (B \times C)$. Therefore $(A - B) \times C \subseteq (A \times C) - (B \times C)$.

(b) Similarly, this inequality is **true**, and we can reverse our steps in part (a) to get a proof.

Let (a, c) be an arbitrary element of $(A \times C) - (B \times C)$. This means that $(a, c) \in A \times C$ but $(a, c) \notin B \times C$. Since $(a, c) \in A \times C$, we know that $a \in A$ and $c \in C$. But since $(a, c) \notin B \times C$ although $c \in C$, we also know $a \notin B$. Thus $a \in A$ and $a \notin B$, which means $a \in A - B$. Thus $(a, c) \in (A - B) \times C$. Therefore $(A \times C) - (B \times C) \subseteq (A - B) \times C$.

(c) Since the inequalities in parts (a) and (b) both hold, we get that the equality in (c) holds for all sets A, B, C .

(d) This inequality is **false** no matter what sets we choose for A and B ! To see this, let A and B be any sets. Notice that the empty set $\emptyset \subseteq A - B$ regardless of what A and B are, so $\emptyset \in \mathcal{P}(A - B)$. However, since $\emptyset \in \mathcal{P}(A)$ and $\emptyset \in \mathcal{P}(B)$, we get $\emptyset \notin \mathcal{P}(A) - \mathcal{P}(B)$. Therefore $\mathcal{P}(A - B) \not\subseteq \mathcal{P}(A) - \mathcal{P}(B)$.

Note. You can prove that if X is any *nonempty* set so that $X \in \mathcal{P}(A - B)$, then $X \in \mathcal{P}(A) - \mathcal{P}(B)$. So the only counterexample to the inequality in part (d) is the empty set.

(e) This inequality is also **false**, but counterexamples are easier to find. For example, let $A = \{1, 2\}$ and $B = \{1\}$. Then $\{1, 2\} \subseteq A$ and $\{1, 2\} \not\subseteq B$, so $\{1, 2\} \in \mathcal{P}(A)$ and $\{1, 2\} \notin \mathcal{P}(B)$, so $\{1, 2\} \in \mathcal{P}(A) - \mathcal{P}(B)$. However $A - B = \{2\}$, so $\{1, 2\} \notin \mathcal{P}(A - B)$. Therefore $\mathcal{P}(A) - \mathcal{P}(B) \not\subseteq \mathcal{P}(A - B)$.

(f) Since the inequality in (e) (or (d)) fails, the equality in (f) fails too.

1. For each positive integer n , let $[n] = \{1, 2, 3, \dots, n\}$, and define

- $\mathcal{S}_\cup(n)$ = the set of all ordered pairs (A, B) of sets such that $A \cup B = [n]$;
- $\mathcal{S}_\cap(n)$ = the set of all ordered pairs (A, B) of subsets of $[n]$ such that $A \cap B = \emptyset$;
- $\mathcal{S}_\subseteq(n)$ = the set of all ordered pairs (A, B) of subsets of $[n]$ such that $A \subseteq B$.

- (a) Find $\mathcal{S}_\cup(1)$ and $\mathcal{S}_\cup(2)$.
- (b) Prove that $\mathcal{S}_\cup(n)$ has exactly 3^n elements.
- (c) Prove that $(A, B) \in \mathcal{S}_\cup(n)$ if and only if $(A^c, B^c) \in \mathcal{S}_\cap(n)$ (here $[n]$ is the universal set). Therefore find the number of elements in $\mathcal{S}_\cap(n)$.
- (d) Prove that $(A, B) \in \mathcal{S}_\cup(n)$ if and only if $(A^c, B) \in \mathcal{S}_\subseteq(n)$ (here $[n]$ is the universal set). Therefore find the number of elements in $\mathcal{S}_\subseteq(n)$.

(a) We get

$$\mathcal{S}_\cup(1) = \{(\{1\}, \emptyset), (\emptyset, \{1\}), (\{1\}, \{1\})\}$$

and

$$\begin{aligned} \mathcal{S}_\cup(2) = \{ & (\{1, 2\}, \emptyset), (\emptyset, \{1, 2\}), (\{1, 2\}, \{1\}), (\{1\}, \{1, 2\}), (\{1, 2\}, \{2\}), \\ & (\{2\}, \{1, 2\}), (\{1, 2\}, \{1, 2\}), (\{1\}, \{2\}), (\{2\}, \{1\}) \}. \end{aligned}$$

- (b) We count how many ways there are to construct sets A and B so that $A \cup B = \{1, 2, \dots, n\}$. To get this union, we need each number from 1 to n to either be in A , or in B , or in both. So we have three possibilities for each of the n numbers from 1 to n . Since these choices are all independent, there are $3 \cdot 3 \cdot \dots \cdot 3 = 3^n$ such ordered pairs (A, B) .
- (c) First assume that $(A, B) \in \mathcal{S}_\cup(n)$. Then $A \cup B = [n]$, so by De Morgan's Law (page 272, #9(a)),

$$A^c \cap B^c = (A \cup B)^c = [n]^c = \emptyset,$$

therefore $(A^c, B^c) \in \mathcal{S}_\cap(n)$.

Conversely, assume that $(A^c, B^c) \in \mathcal{S}_\cap(n)$. Then $A^c \cap B^c = \emptyset$, so by various properties on page 272,

$$A \cup B = (A^c)^c \cup (B^c)^c = (A^c \cap B^c)^c = \emptyset^c = [n],$$

therefore $(A, B) \in \mathcal{S}_\cup(n)$.

This means that there is a one-to-one correspondence between the elements of $\mathcal{S}_\cup(n)$ and the elements of $\mathcal{S}_\cap(n)$, so by part (b) $\mathcal{S}_\cap(n)$ must also have 3^n elements.

- (d) First assume that $(A, B) \in \mathcal{S}_\cup(n)$, which means $A \cup B = [n]$. We want to prove that $(A^c, B) \in \mathcal{S}_\subseteq(n)$, which means we want to prove that $A^c \subseteq B$. Let $x \in A^c$ be arbitrary. This means that $x \in [n]$ but $x \notin A$. Since $A \cup B = [n]$, $x \in [n]$ means $x \in A \cup B$, and since $x \notin A$ we conclude that $x \in B$. Therefore $A^c \subseteq B$ and $(A^c, B) \in \mathcal{S}_\subseteq(n)$.

Conversely, assume that $(A^c, B) \in \mathcal{S}_{\subseteq}(n)$, which means $A^c \subseteq B$. We want to prove that $(A, B) \in \mathcal{S}_{\cup}(n)$, which means we want to prove that $A \cup B = [n]$. Since $A \cup B \subseteq [n]$, we only need to prove that $[n] \subseteq A \cup B$. Let $x \in [n]$ be arbitrary. If $x \in A$, then $x \in A \cup B$ which is what we want. On the other hand, if $x \notin A$, then $x \in A^c$, and since $A^c \subseteq B$, this means that $x \in B$ and thus $x \in A \cup B$. So in either case we get that $x \in A \cup B$. Therefore $[n] \subseteq A \cup B$, so $A \cup B = [n]$, so $(A, B) \in \mathcal{S}_{\cup}(n)$.

Once again this means that there is a one-to-one correspondence between the elements of $\mathcal{S}_{\cup}(n)$ and the elements of $\mathcal{S}_{\subseteq}(n)$, so by part (b) $\mathcal{S}_{\subseteq}(n)$ must also have 3^n elements.

2. For each positive integer n , let $f(n)$ be the number of ordered pairs (A, B) of subsets of $\{1, 2, 3, \dots, n\}$ so that $A \cup B$ has an even number of elements.

- (a) Find $f(1)$ and $f(2)$ by listing all the ordered pairs of subsets.
 (b) Use Problem 1(b) to prove that for any n ,

$$f(n) = \sum_{k=0}^{\lfloor n/2 \rfloor} \binom{n}{2k} 3^{2k}.$$

Show that your answers to part (a) agree with this formula.

- (c) Mimic Example 6.7.4 on page 368 to prove that $\sum_{i=0}^n \binom{n}{i} 3^i = 4^n$ and thus

$$\sum_{k=1}^{\lfloor (n+1)/2 \rfloor} \binom{n}{2k-1} 3^{2k-1} = 4^n - f(n).$$

- (d) Use Pascal's Formula (page 360), (b) and (c), and mathematical induction to prove that

$$f(n) = \begin{cases} 2^{n-1}(2^n - 1) & \text{if } n \text{ is odd,} \\ 2^{n-1}(2^n + 1) & \text{if } n \text{ is even.} \end{cases}$$

- (a) Since A and B are subsets of $\{1, 2, \dots, n\}$, we always have $A \cup B \subseteq \{1, 2, \dots, n\}$. So when $n = 1$, the only way for $A \cup B$ to have an even number of elements is if $A \cup B = \emptyset$, so the only ordered pair (A, B) that works is (\emptyset, \emptyset) , and thus $f(1) = \mathbf{1}$. When $n = 2$, we could have $A \cup B = \emptyset$ or $A \cup B = \{1, 2\}$, so the ordered pairs (A, B) that work are (\emptyset, \emptyset) plus the nine ordered pairs in $\mathcal{S}_{\cup}(2)$ from problem 1(a). Thus $f(2) = \mathbf{10}$.
- (b) First, from problem 1(b) it is clear that for any set S with m elements there must be exactly 3^m ordered pairs (A, B) of sets so that $A \cup B = S$ (since the names of the m elements of S don't matter). Let k be an integer so that $0 \leq 2k \leq n$. There are $\binom{n}{2k}$ subsets of $\{1, 2, \dots, n\}$ with $2k$ elements, and for each of these subsets there are 3^{2k} ordered pairs (A, B) of sets whose union is that subset. Thus for each k , there are $\binom{n}{2k} 3^{2k}$ ordered pairs (A, B) of subsets of $\{1, 2, \dots, n\}$ so that $A \cup B$ has $2k$ elements. Adding over all possible values of k (namely $k = 0, 1, \dots, \lfloor n/2 \rfloor$), we get that

$$f(n) = \sum_{k=0}^{\lfloor n/2 \rfloor} \binom{n}{2k} 3^{2k}.$$

When $n = 1$ this says

$$f(1) = \sum_{k=0}^0 \binom{1}{2k} 3^{2k} = \binom{1}{0} 3^0 = 1,$$

and when $n = 2$ it says

$$f(2) = \sum_{k=0}^1 \binom{2}{2k} 3^{2k} = \binom{2}{0} 3^0 + \binom{2}{2} 3^2 = 1 + 9 = 10,$$

both agreeing with part (a).

- (c) We put $a = 1$ and $b = 3$ into the Binomial Theorem (Theorem 6.7.1 on page 364) to get

$$\sum_{i=0}^n \binom{n}{i} 3^i = \sum_{i=0}^n \binom{n}{i} 1^{n-i} 3^i = (1 + 3)^n = 4^n.$$

Splitting this sum into two parts, one with all the even i 's and one with all the odd i 's, we get

$$\sum_{k=0}^{\lfloor n/2 \rfloor} \binom{n}{2k} 3^{2k} + \sum_{k=1}^{\lfloor (n+1)/2 \rfloor} \binom{n}{2k-1} 3^{2k-1} = 4^n.$$

But the first sum is just $f(n)$ by part (b), so subtracting it from both sides gives us

$$\sum_{k=1}^{\lfloor (n+1)/2 \rfloor} \binom{n}{2k-1} 3^{2k-1} = 4^n - f(n)$$

as required.

- (d) *Basis step.* When $n = 1$ (which is odd) the formula says $f(1) = 2^0(2^1 - 1) = 1$, which is correct by part (a).

Inductive step. Assume that the formula is correct for some integer $n \geq 1$. We want to prove it is correct for the next integer $n + 1$. Well,

$$\begin{aligned} f(n+1) &= \sum_{k=0}^{\lfloor (n+1)/2 \rfloor} \binom{n+1}{2k} 3^{2k} && \text{by part (b)} \\ &= \sum_{k=0}^{\lfloor (n+1)/2 \rfloor} \left[\binom{n}{2k} + \binom{n}{2k-1} \right] 3^{2k} && \text{by Pascal's Formula} \\ &= \sum_{k=0}^{\lfloor n/2 \rfloor} \binom{n}{2k} 3^{2k} + \sum_{k=1}^{\lfloor (n+1)/2 \rfloor} \binom{n}{2k-1} 3 \cdot 3^{2k-1} \\ &= f(n) + 3(4^n - f(n)) && \text{by parts (b) and (c)} \\ &= 3(4^n) - 2f(n) \\ &= \begin{cases} 3(4^n) - 2^n(2^n - 1) & \text{if } n \text{ is odd} \\ 3(4^n) - 2^n(2^n + 1) & \text{if } n \text{ is even} \end{cases} && \text{by assumption} \\ &= \begin{cases} 2(4^n) + 2^n = 2^n(2^{n+1} + 1) & \text{if } n+1 \text{ is even} \\ 2(4^n) - 2^n = 2^n(2^{n+1} - 1) & \text{if } n+1 \text{ is odd,} \end{cases} \end{aligned}$$

which completes the inductive step. Therefore the formula is correct for all integers $n \geq 1$.

Note: If n is odd, and if $2^n - 1$ happens to be a prime number, then the value $f(n) = 2^{n-1}(2^n - 1)$ is what is called a *perfect number*. To find out what these are, ask your professor or TA, or search the internet.

3. Again let $[n] = \{1, 2, 3, \dots, n\}$ for any positive integer n .

- (a) Find the number of functions $f : [n] \rightarrow [n]$ such that $f(k) \leq k \forall k \in [n]$.
 - (b) Find the number of one-to-one functions $f : [n] \rightarrow [n]$ such that $f(k) \leq k \forall k \in [n]$.
 - (c) Find the number of functions $f : [n] \rightarrow [n]$ such that $f(k) \leq k + 1 \forall k \in [n]$.
 - (d) Find the number of onto functions $f : [n] \rightarrow [n]$ such that $f(k) \leq k + 1 \forall k \in [n]$.
- (a) Since, for every k , $f(k)$ must be one of the k values $1, 2, \dots, k$, there is one choice for $f(1)$ (namely 1), two choices for $f(2)$ (namely 1 or 2), and so on up to n choices for $f(n)$ (namely any of $1, 2, \dots, n$). Thus by the Multiplication Rule there are $1 \cdot 2 \cdot \dots \cdot n = n!$ ways to assign all the values $f(1), f(2), \dots, f(n)$, that is, **$n!$** different functions.
 - (b) If f must be one-to-one, then we still must assign $f(1) = 1$, but then we cannot assign $f(2)$ to be 1 too, so we must put $f(2) = 2$. Next we cannot let $f(3)$ be 1 or 2, so we must put $f(3) = 3$. Continuing in this way, we are forced to put $f(k) = k$ for each k , so there is just **one** one-to-one function $f : [n] \rightarrow [n]$, namely the identity function.
 - (c) Proceeding as in part (a), for each k , $f(k)$ must be one of the $k+1$ choices $1, 2, \dots, k+1$, provided that $k < n$. So $f(1)$ can be 1 or 2, $f(2)$ can be 1, 2 or 3, and so on up to $f(n-1)$ which can be any of $1, 2, \dots, n$. But $f(n)$ must still belong to $[n]$ so there are only n choices for $f(n)$. Thus by the Multiplication Rule the total number of functions is $2 \cdot 3 \cdot \dots \cdot n \cdot n = \mathbf{n(n!)}$.
 - (d) Note that since $[n]$ is finite, a function $f : [n] \rightarrow [n]$ is onto if and only if it is one-to-one. So we are really just counting one-to-one functions again. Now $f(1)$ must be 1 or 2, so there are two choices for $f(1)$. Then $f(2)$ must be 1, 2 or 3, so removing whichever choice we made for $f(1)$ will leave two choices for $f(2)$. In general there will be $k+1$ choices for $f(k)$ (namely $1, 2, \dots, k+1$), but after we remove the choices we make for $f(1), f(2), \dots, f(k-1)$ we will always have exactly two choices left for $f(k+1)$. The exception again is that for $f(n)$ there are only n choices originally (namely $1, 2, \dots, n$), and after we remove the choices we make for $f(1), f(2), \dots, f(n-1)$ we will only have one choice left for $f(n)$. So in total there will be $2 \cdot 2 \cdot \dots \cdot 2 \cdot 1 = \mathbf{2^{n-1}}$ onto functions.

1. If $f : \mathbf{R} \rightarrow \mathbf{R}$ is a function (where \mathbf{R} is the set of all real numbers), we define the function $f^{(2)}$ to be the composition $f \circ f$, and for any integer $n \geq 2$, define $f^{(n+1)} = f \circ f^{(n)}$. So $f^{(2)}(x) = (f \circ f)(x) = f(f(x))$, $f^{(3)}(x) = (f \circ f^{(2)})(x) = f(f(f(x)))$, and so on.

Let $f : \mathbf{R} \rightarrow \mathbf{R}$ be defined by $f(x) = 3x^2$ for all $x \in \mathbf{R}$.

- Find and simplify $f^{(2)}(x)$ and $f^{(3)}(x)$.
- Use part (a) (and more calculations if you need them) to guess a formula for $f^{(n)}(x)$.
- Prove your guess using mathematical induction.
- Find all $x \in \mathbf{R}$ so that $f^{(271)}(x) = x$.

- (a) We get

$$f^{(2)}(x) = f(f(x)) = f(3x^2) = 3(3x^2)^2 = 3^3x^4$$

and

$$f^{(3)}(x) = f(f^{(2)}(x)) = f(3^3x^4) = 3(3^3x^4)^2 = 3^7x^8.$$

- (b) Since $f^{(2)}(x) = 3^3x^4 = 3^{2^2-1}x^{2^2}$ and $f^{(3)}(x) = 3^7x^8 = 3^{2^3-1}x^{2^3}$, we guess that

$$f^{(n)}(x) = 3^{2^n-1}x^{2^n} \quad \text{for all } n \geq 2.$$

- (c) *Basis step.* This is already done, since our formula is true when $n = 2$.

Inductive step. Assume that $f^{(k)}(x) = 3^{2^k-1}x^{2^k}$ for some integer $k \geq 2$. We want to prove that $f^{(k+1)}(x) = 3^{2^{k+1}-1}x^{2^{k+1}}$. Well,

$$\begin{aligned} f^{(k+1)}(x) &= f\left(f^{(k)}(x)\right) \quad \text{by definition} \\ &= f\left(3^{2^k-1}x^{2^k}\right) \quad \text{by assumption} \\ &= 3\left(3^{2^k-1}x^{2^k}\right)^2 \\ &= 3^{1+(2^k-1)2}x^{(2^k)2} \\ &= 3^{2^{k+1}-1}x^{2^{k+1}}, \end{aligned}$$

which finishes the inductive step. This proves that our guess is correct for every integer $n \geq 2$.

- (d) Since $f^{(271)}(x) = 3^{2^{271}-1}x^{2^{271}}$, we need to solve the equation $3^{2^{271}-1}x^{2^{271}} = x$. One obvious solution is $x = \mathbf{0}$. So assuming now that $x \neq 0$, we can divide both sides by x and get $3^{2^{271}-1}x^{2^{271}-1} = 1$, which can be rewritten as $(3x)^{2^{271}-1} = 1$, which means that $3x = 1$ since $2^{271} - 1$ is odd. Thus the only other solution is $x = \mathbf{1/3}$.

2. Let $[n] = \{1, 2, \dots, n\}$, where n is a positive integer. Let \mathcal{R} be the relation on the power set $\mathcal{P}([n])$ defined by: for $A, B \in \mathcal{P}([n])$, $A\mathcal{R}B$ if and only if $1 \notin A - B$.

- (a) Is \mathcal{R} reflexive? Symmetric? Transitive? Explain.

- (b) Find the number of ordered pairs (A, B) of sets in $\mathcal{P}([n])$ such that ARB . [*Hint*: first count the number of ordered pairs (A, B) of sets in $\mathcal{P}([n])$ so that $A\mathcal{R}B$.]
- (c) Suppose you choose sets $A, B \in \mathcal{P}([n])$ at random. What is the probability that ARB ?
- (d) Let \mathcal{S} be the relation on the power set $\mathcal{P}([n])$ defined by: for $A, B \in \mathcal{P}([n])$, ASB if and only if $1 \in A - B$. Is \mathcal{S} transitive? Explain.

- (a) **\mathcal{R} is reflexive.** Here is a proof. Let $A \in \mathcal{P}([n])$ be arbitrary. Then $A - A = \emptyset$, so $1 \notin A - A$, so ARA .

\mathcal{R} is not symmetric. Here is a counterexample. Let $A = \emptyset$ and $B = \{1\}$. Then $A - B = \emptyset$, so $1 \notin A - B$, so ARB . However $B - A = \{1\}$, so $1 \in B - A$, so $B\mathcal{R}A$.

\mathcal{R} is transitive. Here is a proof. Let $A, B, C \in \mathcal{P}([n])$ be arbitrary so that ARB and BRC . This means that $1 \notin A - B$ and $1 \notin B - C$. We want to prove that ARC , which means we want to prove that $1 \notin A - C$. We do this by contradiction. Suppose that $1 \in A - C$. This means that $1 \in A$ but $1 \notin C$. Since $1 \in A$ but $1 \notin A - B$, it must mean that $1 \in B$. But now $1 \in B$ and $1 \notin B - C$ means $1 \in B - C$, which is a contradiction. Therefore $1 \notin A - C$, so \mathcal{R} is transitive.

- (b) Since $A\mathcal{R}B$ means $1 \in A - B$, to count the number of ordered pairs (A, B) so that $A\mathcal{R}B$ we just count the number of (A, B) so that $1 \in A$ and $1 \notin B$. The number of subsets A of $[n]$ so that $1 \in A$ is just the number of subsets of $\{2, 3, \dots, n\}$, which is 2^{n-1} (for example, see p. 285). The number of subsets B of $[n]$ so that $1 \notin B$ is also just the number of subsets of $\{2, 3, \dots, n\}$, which is 2^{n-1} . Thus the number of ordered pairs (A, B) so that $A\mathcal{R}B$ is $2^{n-1} \cdot 2^{n-1} = 2^{2(n-1)}$ by the Multiplication Rule. There are 2^n subsets of $[n]$ altogether, so there are $2^n \cdot 2^n = 2^{2n}$ ordered pairs (A, B) altogether. Therefore the number of ordered pairs (A, B) of sets in $\mathcal{P}([n])$ such that ARB is $2^{2n} - 2^{2(n-1)} = 2^{2n-2}(2^2 - 1) = 3(2^{2n-2})$.

- (c) Since all choices of subsets $A, B \in \mathcal{P}([n])$ are equally likely, the probability is

$$\frac{\text{number of } (A, B) \text{ so that } ARB}{\text{total number of } (A, B)} = \frac{3(2^{2n-2})}{2^{2n}} = \frac{3}{4},$$

regardless of the value of n .

- (d) Yes, \mathcal{S} is transitive, **vacuously**. Suppose that $A, B, C \in \mathcal{P}([n])$ satisfy ASB and BSC . This means that $1 \in A - B$ and $1 \in B - C$. But $1 \in A - B$ means in particular that $1 \notin B$, while $1 \in B - C$ means in particular that $1 \in B$. This is a contradiction, so the “if” part of the definition of transitivity can never happen, so the relation \mathcal{S} is transitive vacuously.

3. Let \mathcal{F} be the set of all functions $f : \{1, 2, \dots, n\} \rightarrow \{1, 2, \dots, n\}$, where n is a positive integer. Define a relation R on \mathcal{F} by: for $f, g \in \mathcal{F}$, fRg if and only if $f(k) + g(k)$ is even for all $k \in \{1, 2, \dots, n\}$.

- (a) Prove that R is an equivalence relation on \mathcal{F} .
- (b) Suppose that $n = 2m + 1$ is odd. Find the number of functions in the equivalence class $[id]$, where id is the identity function on $\{1, 2, \dots, n\}$. How many of these functions are one-to-one and onto?

(c) Suppose that $n = 2m$ is even. Find the number of functions in the equivalence class $[g]$, where $g(x) = 1$ is a constant function. How many of these functions are one-to-one and onto?

(a) *R is reflexive.* Let $f \in \mathcal{F}$ be arbitrary. Then $f(k) + f(k) = 2f(k)$ is even for every $k \in \{1, 2, \dots, n\}$, since $f(k)$ is an integer, so fRf .

R is symmetric. Let $f, g \in \mathcal{F}$ be arbitrary so that fRg . This means that $f(k) + g(k)$ is even for all $k \in \{1, 2, \dots, n\}$. But then $g(k) + f(k) = f(k) + g(k)$ is even for all $k \in \{1, 2, \dots, n\}$, so gRf .

R is transitive. Let $f, g, h \in \mathcal{F}$ be arbitrary so that fRg and gRh . This means that $f(k) + g(k)$ is even for all $k \in \{1, 2, \dots, n\}$, and $g(k) + h(k)$ is even for all $k \in \{1, 2, \dots, n\}$. But then $f(k) + g(k) + g(k) + h(k) = f(k) + h(k) + 2g(k)$ is even for all $k \in \{1, 2, \dots, n\}$, so $f(k) + h(k)$ is even for all $k \in \{1, 2, \dots, n\}$, since the sum and difference of even numbers is even. Therefore fRh .

(b) We want to count the number of functions $f \in \mathcal{F}$ so that $fRid$. *id* is the function $id(k) = k$ for all $k \in \{1, 2, \dots, n\}$. Thus $fRid$ means that $f(k) + k$ is even for all k . This in turn means that $f(k)$ must be even whenever k is even, and odd whenever k is odd. Since $n = 2m + 1$, there are m even numbers and $m + 1$ odd numbers in $\{1, 2, \dots, n\}$. So for each of the m even k 's there are m choices for $f(k)$, and for each of the $m + 1$ odd k 's there are $m + 1$ choices for $f(k)$. Thus the total number of ways we can define f is $\mathbf{m^m(m + 1)^{m+1}}$.

If we insist that f be one-to-one and onto, we only need to make it one-to-one, since any one-to-one function $f : \{1, 2, \dots, n\} \rightarrow \{1, 2, \dots, n\}$ will have to be onto as well. Now when we count how many ways there are to define $f(2), f(4), \dots, f(2m)$ (that is, $f(k)$ for the m even k 's), we get m ways to define $f(2)$, $m - 1$ ways to define $f(4)$, and so on down to just one way to define $f(2m)$. So there are $m!$ ways to define $f(2), f(4), \dots, f(2m)$. Similarly, there are $m + 1$ ways to define $f(1)$, m ways to define $f(3)$, and so on down to just one way to define $f(2m + 1)$. So there are $(m + 1)!$ ways to define $f(1), f(3), \dots, f(2m + 1)$. Thus altogether there are $\mathbf{m!(m + 1)!}$ ways to define f so that it is one-to-one and onto.

(c) This time we want to count the number of functions $f \in \mathcal{F}$ so that fRg . But since $g(k) = 1$ for all $k \in \{1, 2, \dots, n\}$, to get $f(k) + g(k)$ to be even for all k , we will need that $f(k)$ is odd for all k . Since $n = 2m$, there are m odd numbers in $\{1, 2, \dots, n\}$. So we have m choices for each $f(k)$, and thus the total number of ways we can define f is $\mathbf{m^n = m^{2m}}$.

If we insist that f be one-to-one and onto, once again we only need to make it one-to-one. But since $f(k)$ must be odd for every k , and the total number of k 's ($2m$) is greater than the number of odd numbers available (m), it is impossible to assign a different odd number to each $f(k)$. Thus the number of one-to-one onto functions this time is **zero**.